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PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP -- A ZERO-ONE I--ETC(U)  
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SERIAL T-423 NL

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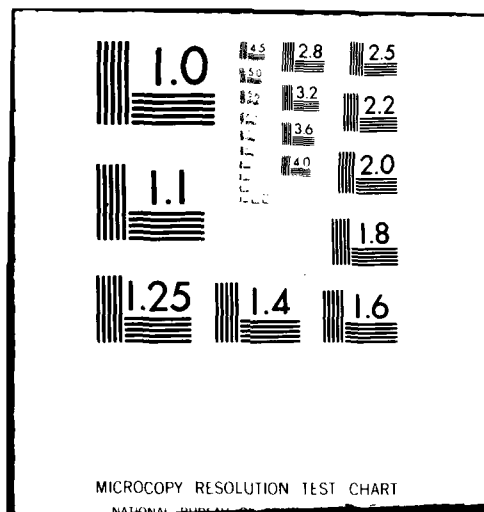
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PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP --  
A ZERO-ONE INTEGER PROGRAM TO SOLVE MULTIACTIVITY  
MULTIFACILITY CAPACITY-CONSTRAINED ASSIGNMENT  
PROBLEMS.

by

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Krishan L. Chhabra  
Richard M. Soland

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Serial T-423

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The George Washington University  
School of Engineering and Applied Science  
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Program in Logistics

Contract N00014-75-C-0729

Project NR 347 020

Office of Naval Research

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SECURITY CLASSIFICATION OF THIS PAGE (When Data Entered)

REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER T-423 ✓	2. GOVT ACCESSION NO. AD-A091	3. RECIPIENT'S CATALOG NUMBER 127
4. TITLE (and Subtitle) PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP-- A ZERO-ONE INTEGER PROGRAM TO SOLVE MULTIACTIVITY MULTIFACILITY CAPACITY-CONSTRAINED ASSIGNMENT PROBLEMS		5. TYPE OF REPORT & PERIOD COVERED SCIENTIFIC
7. AUTHOR(s) KRISHAN L. CHHABRA RICHARD M. SOLAND		6. PERFORMING ORG. REPORT NUMBER T-423
9. PERFORMING ORGANIZATION NAME AND ADDRESS THE GEORGE WASHINGTON UNIVERSITY PROGRAM IN LOGISTICS WASHINGTON, DC 20037		8. CONTRACT OR GRANT NUMBER(s) N00014-75-C-0729 ✓
11. CONTROLLING OFFICE NAME AND ADDRESS OFFICE OF NAVAL RESEARCH CODE 434 ARLINGTON, VIRGINIA 22217		10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS
14. MONITORING AGENCY NAME & ADDRESS (if different from Controlling Office)		12. REPORT DATE 1 July 1980
		13. NUMBER OF PAGES 66
		15. SECURITY CLASS. (of this report) NONE
		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE
16. DISTRIBUTION STATEMENT (of this Report)  APPROVED FOR PUBLIC SALE AND RELEASE: DISTRIBUTION UNLIMITED.		
17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report)		
18. SUPPLEMENTARY NOTES		
19. KEY WORDS (Continue on reverse side if necessary and identify by block number) INTEGER PROGRAMMING MULTIACTIVITY MULTIFACILITY ASSIGNMENT PROBLEM		
20. ABSTRACT (Continue on reverse side if necessary and identify by block number) ZIPCAP is a zero-one integer program developed to solve multiactivity multifacility capacity-constrained assignment problems. The objective in such problems is an assignment that minimizes the sum of the variable costs and fixed costs. This document describes ZIPCAP and provides guidelines for its use. Illustrative areas of application, program options, and test problems are included.		

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S/N 0102-014-6601

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SECURITY CLASSIFICATION OF THIS PAGE (When Data Entered)

THE GEORGE WASHINGTON UNIVERSITY  
School of Engineering and Applied Science  
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Program in Logistics

Abstract  
of  
Serial T-423  
1 July 1980

PROGRAM DESCRIPTION AND USER'S GUIDE FOR ZIPCAP --  
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MULTIFACILITY CAPACITY-CONSTRAINED ASSIGNMENT  
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ZIPCAP is a zero-one integer program developed to solve multi-activity multifacility capacity-constrained assignment problems. The objective in such problems is an assignment that minimizes the sum of variable costs and fixed costs. This document describes ZIPCAP and provides guidelines for its use. Illustrative areas of application, program options, and test problems are included.

Research Supported by

Contract N00014-75-C-0729  
Project NR 347 020  
Office of Naval Research

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1. Introduction

Multiactivity multifacility assignment problems are important because of their practical applications. The objective of such a problem is to minimize the total cost of the system, which includes both variable costs and fixed costs. A problem which has been of recent interest takes into account capacity constraints as well [Gross, Pinkus, and Soland (1979)]. A solution algorithm and a computer program implementing it have been developed for this kind of problem, i.e., for multiactivity multifacility 0-1 assignment problems with capacity constraints.

The development of the solution algorithm and computational test results using the computer program are described in a separate report titled "Solving a Multiactivity Multifacility Capacity-constrained 0-1 Assignment Problem."

This document, on the other hand, provides a description of the

computer program and instructions for its use. It is a FORTRAN program and is named ZIPCAP, an acronym for Zero-one Integer Program to solve multi-activity multifacility Capacity-constrained Assignment Problems.

Section 2 of this document provides a mathematical formulation of the basic problem and examples of application areas. Section 3 includes the computational steps and a program description, whereas Section 4 provides user information including program options and input data specifications. Test problems and output illustrations are covered in Section 5.

## 2. Problem Formulation

### 2.1 Basic Problem

The basic problem, called problem (P), is to find  $x_{ij}$  and  $y_k$  values that:

$$(P) \left\{ \begin{array}{ll} \text{Minimize} & \sum_{i=1}^m \sum_{j=1}^n a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k \quad (1) \\ \text{subject to} & \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n \quad (2) \\ & \sum_{i=1}^m \sum_{j=1}^n d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p \quad (3) \\ & x_{ij}, y_k = 0 \text{ or } 1 \text{ for all } i, j, k \quad (4) \end{array} \right.$$

where  $m$ ,  $n$ , and  $p$  represent the number of designs, activities, and facilities respectively; and  $i$ ,  $j$ ,  $k$  are the corresponding indices.

The parameters are interpreted as follows:

$a_{ij}$  = variable cost of activity  $j$  using design  $i$  ,

$b_k$  = fixed cost of facility  $k$  ,

$s_k$  = capacity available at facility  $k$  , and

$d_{ijk}$  = capacity required at facility  $k$  for activity  $j$   
when activity  $j$  uses design  $i$  .

The decision variables are:

$x_{ij}$  = 1 if activity  $j$  uses design  $i$  ,  
= 0 otherwise.

$y_k$  = 1 if facility  $k$  is used ,  
= 0 otherwise.

Simply stated, in problem (P) one seeks to minimize the sum of variable and fixed costs subject to constraint set (2) that each activity be assigned a single design and subject to constraint set (3) that the capacity constraints imposed by the facility capacities are satisfied.

Problem (P) has  $mn+p$  0-1 variables and  $n+p$  constraints.

## 2.2 Key Elements

The key elements in problem (P) are activities, facilities, designs, variable costs, and fixed costs. In general, these can be defined as follows.

- . Activity -- a type of component, product or item under consideration.

- . Facility -- an installation or location where activities can be served or serviced.
- . Design -- a meaningful configuration of facilities along with a meaningful strategy for using them. For example, if there are four facilities 1, 2, 3 and 4, a design may include facilities 1, 2 and 4, whereas another design may include facility 2 only.
- . Variable Cost -- the cost associated with a particular activity being assigned to a particular design.
- . Fixed Cost -- the cost associated with a facility (such as part of its purchase, operation and maintenance cost) that is independent of the activities served at that facility.

### 2.3 Areas of Application

The computer program ZIPCAP is designed to solve a problem that can be formulated as problem (P). The formulation applies to existing as well as to proposed facilities. In other words, it is useful for a situation where the decision may be to delete some of the existing facilities, as well as for a situation involving a choice among proposed facilities (or locations).

Table 1 includes examples of areas where this formulation can be applied. Within each application area, activities and facilities are identified. The implications of designs, variable costs and fixed costs are apparent.

### 3. Computational Procedure and Program Description

The procedure to solve problem (P) basically involves branch and bound [Cooffrion and Marsten (1972)], and Lagrangean relaxation

Table 1

EXAMPLES OF APPLICATION AREAS

Area of Application	Activities	Facilities
1. Design of multi-echelon inventory systems <sup>*</sup>	Items to be stocked	Warehouses (Comprising different levels or echelons, e.g., central, regional, and local warehouses)
2. Assignment of repairable components <sup>**</sup>	Major components of a unit, e.g., components of an aircraft, a ship, a piece of machinery	Repair depots
3. Design of training programs	Training program categories or occupational classifications	Training schools
4. Location of facilities <sup>***</sup>	Types of services, e.g., health-care services	Buildings or installations, e.g., health-care centers

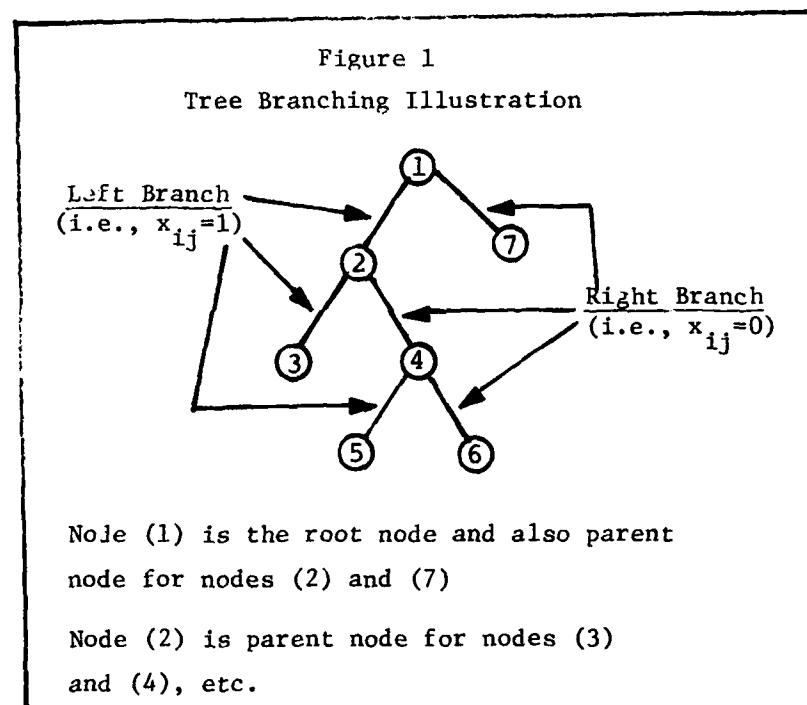
<sup>\*</sup>Gross, Pinkus, and Soland (1979)

<sup>\*\*</sup>Gross and Pinkus (1979)

<sup>\*\*\*</sup>Pinkus, Gross, and Soland (1973)

[Geoffrion (1974)]. In addition certain heuristics are used for branching, and to exclude infeasible assignments.

Branching is done on the  $x_{ij}$  variables by setting them equal to 1 or 0. The branching commences by setting a variable equal to 1 and moving to the left-branch node. When backtracking, the corresponding variable is set equal to 0 and we move to the right-branch node (if the right-branch has not yet been explored). Figure 1 illustrates a branching tree.



The  $x_{ij}$  variables which have been set equal to 1 or 0 at any given node are termed fixed variables, and the remaining ones are termed free variables.

### 3.1 Computational Steps

Figure 2 presents a simplified flow diagram, showing the

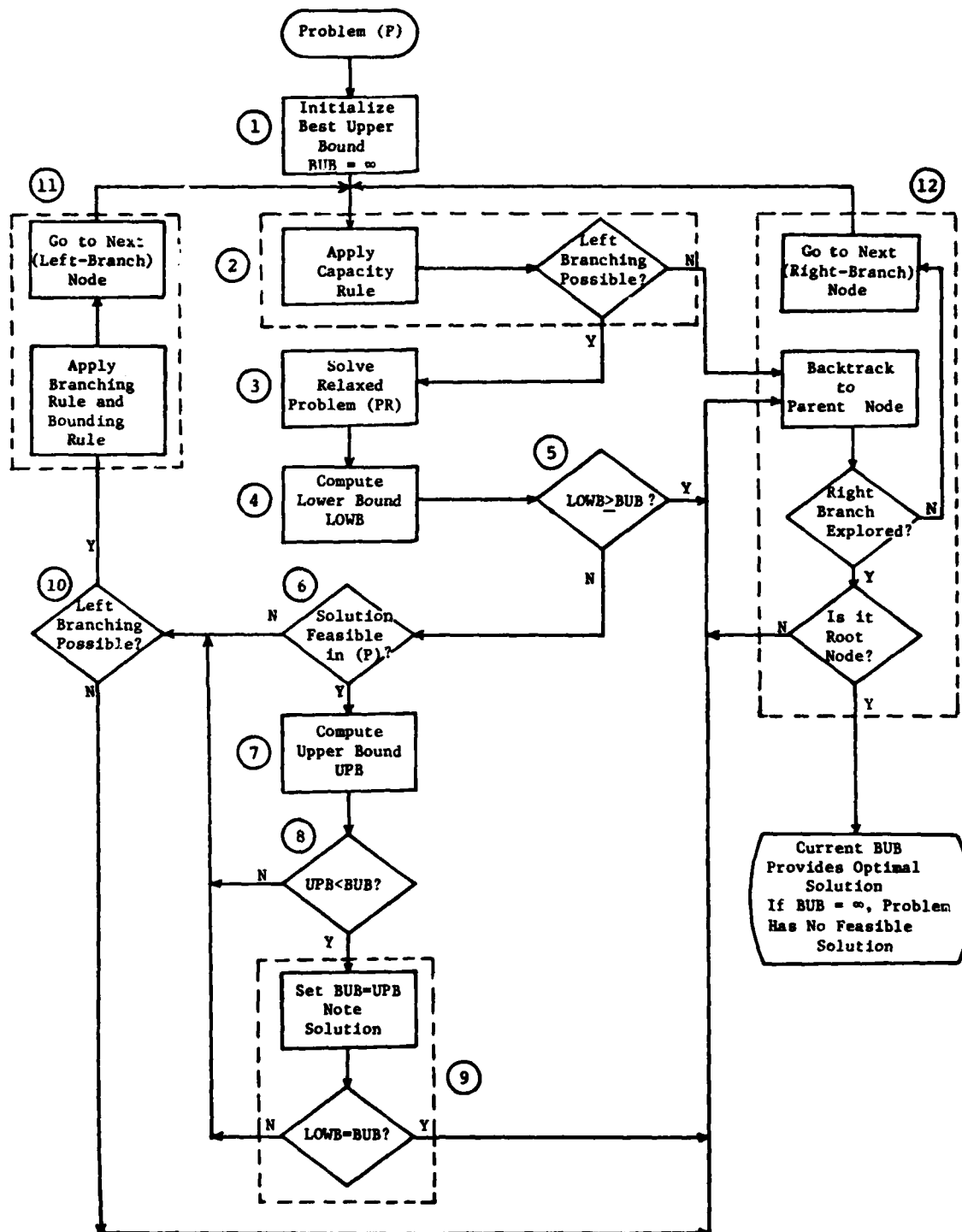


Figure 2  
Simplified Flow Diagram Showing Computational Steps

computational steps of the computer program. Following is a description of these steps.

Step 1. Initialize by setting best upper bound (BUB) equal to a large value and the node number to 1.

Step 2. Apply the capacity rule to exclude infeasible assignments prior to solving the relaxed problem.

Exclude a free  $x_{ij}$  if, for any facility  $k$  and activity  $j$ ,

$$\left( d_{ijk} - \min_j d_{ijk} \right) > \left( s_k - \sum_j \min_j d_{ijk} \right).$$

If an  $x_{ij}$  is already fixed as 1, the corresponding  $d_{ijk}$  replaces  $\min_j d_{ijk}$ .

If  $\sum_j \min_j d_{ijk} > s_k$ , then backtrack, i.e., go to step 12.

Step 3. Solve the relaxed problem (PR), i.e.,

$$(PR) \begin{cases} \text{Minimize} & \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \\ \text{subject to} & \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n, \\ & x_{ij} = 0 \text{ or } 1 \text{ for all } i, j, \end{cases}$$

$$\text{where } c_{ij} = a_{ij} + \sum_{k=1}^p b_k \left( 1 - \delta_k \right) \left( d_{ijk} / s_k \right),$$

$$\delta_k = 1 \text{ if } \sum_i \sum_j x_{ij} e_{ik} > C,$$

such that  
 $x_{ij}$  is fixed

= 0 otherwise, and



$$e_{ik} = 1 \text{ if design } i \text{ uses facility } k ,$$

$$= 0 \text{ otherwise.}$$

Step 4. Lower bound (LOWB) is the sum of the optimal solution value of problem (PR) and the fixed cost FC. FC is the cost of the facilities forced into the solution because of the  $x_{ij}$  variables fixed as 1:

$$FC = \sum_{k=1}^p \delta_k b_k$$

Step 5. Compare lower bound with best upper bound. If  $LOWB \geq BUB$ , go to step 12.

Step 6. Check if the solution of problem (PR) obtained in step 3 satisfies the capacity constraints of problem (P), i.e.,

$$\sum_{i=1}^m \sum_{j=1}^n d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p,$$

$$\text{where } y_k = 1 \text{ if } \sum_i \sum_j d_{ijk} x_{ij} > 0 ,$$

$$= 0 \text{ otherwise}$$

If the capacity constraints are not satisfied, go to step 10.

Step 7. Compute upper bound (UPB), i.e.,

$$UPB = \sum_{i=1}^m \sum_{j=1}^n a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k$$

Step 8. Compare upper bound with best upper bound. If  $UPB \geq BUB$ , go to step 10.

Step 9. Retain  $UPB$  as  $BUB$ , i.e., set  $BUB = UPB$  and note the corresponding solution comprising  $x_{ij}$  and  $y_k$  variables.

If  $LOWB$  equals  $BUB$  (same as  $UPB$  in this case), go to step 12.

Step 10 If an  $x_{ij}$  variable has been fixed as 1 for each of the  $n$  activities, then go to step 12.

Step 11. Select an  $x_{ij}$  variable for branching to the left node. The choice of the branching variable depends on the  $c_{ij}$  values and is such that the corresponding  $x_{ij}$ , if perturbed, has the maximum impact on the optimal value of problem (PR). That is, for each  $j$ , obtain  $D_j = c_{2j} - c_{1j}$  where  $c_{2j}$  is the second smallest permissible and  $c_{1j}$  is the smallest permissible  $c_{ij}$  value in column  $j$ .

A  $c_{ij}$  value is considered permissible if it does not correspond to an activity  $j$  having  $x_{ij}$  variable fixed as 1, or to an  $x_{ij}$  variable fixed as 0. The  $x_{ij}$  variables can be fixed as 1 or 0 because of the branching rule, backtracking, and/or use of the capacity rule. Further,  $x_{ij}$  may be fixed as 0 because of the bounding rule, i.e., if  $(c_{ij} - c_{1j}) > (BUB - LOWB)$ . The  $x_{ij}$  variable selected has the maximum value of  $D_j$  and has  $c_{ij} = c_{1j}$ .

Set the selected  $x_{ij}$  variable to 1. Increase the node number by one, and go to step 2.

Step 12. Backtracking continues until a node is reached for which

the right branch has not been explored. Set the corresponding  $x_{ij}$  variable to 0. Increase the node number by one, and go to step 2.

If no such node is found, it implies that the root node has been reached, and the procedure terminates.

The current BUB and the solution obtained in step 9 provide the optimal solution and its value. If BUB happens to be the initial value, the problem does not have a feasible solution, i.e., the capacity constraints cannot be satisfied.

### 3.2 Computer Program

The computer program ZIPCAP is written in FORTRAN IV. It is based on the flow diagram and the computational steps described in the preceding paragraphs.

Appendix A provides a listing of the program. Extensive use of comment cards provides details within each computational step.

A dictionary of the symbolic names used in the program listing is provided in Appendix B.

The program has been developed and tested on an IBM 3031 at The George Washington University.\* The only internally supplied subroutine used by the program is TIMET from the PII Library. This subroutine provides the elapsed time for program execution and excludes the time to read and write the input and output.

---

\* IBM 3031 CPU with 3 million Bytes Real Memory. Operating System OS/VS1 Release 6.7. Various compilers including FORTRAN levels G 1 to 6 and H 1 to 6.

The program comprising about 480 lines is currently dimensioned for a problem size of 35 designs (m), 35 activities (n) and 30 facilities (p). The user capacity required to execute this program is 346 K bytes. The functional relationship between the dimensions of the arrays and the storage requirement is given by

$$f(m, n, p) = 4 \left[ (p+4)mn + (m+5)p + 9n \right] \text{ bytes.}$$

Thus the program size to execute a problem has two components: one, due to the program itself, comprises 173 K bytes, and the other, dependent on the dimensions of the arrays, is given by the above functional relationship.

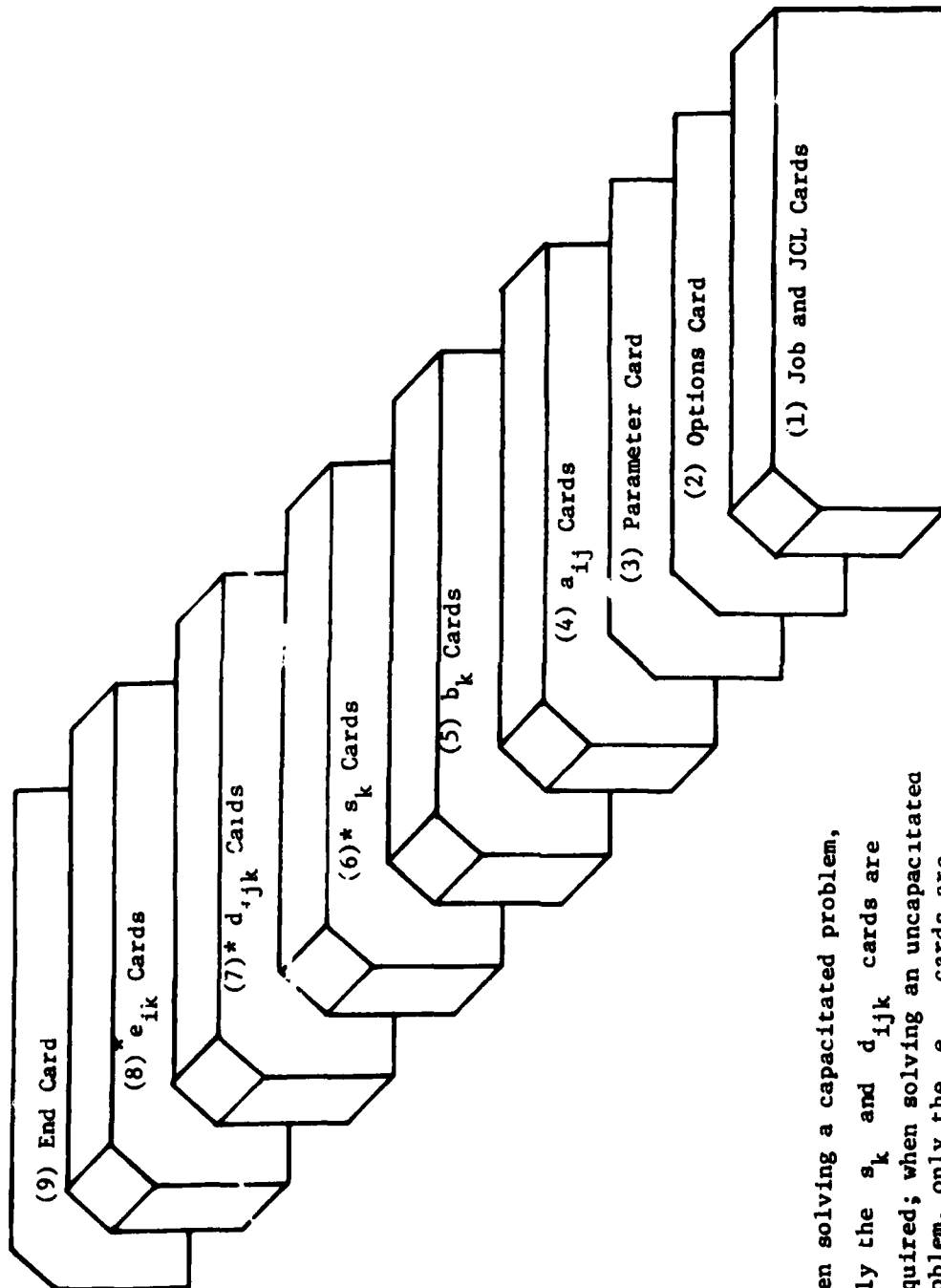
#### 4. User Information

ZIPCAP is designed to solve a multiactivity multifacility capacity-constrained 0-1 assignment problem having formulation (P). Such a problem is referred to as "capacitated" because the optimal solution must satisfy the capacity constraints. However, the program can be used in a situation where the capacity constraints are not applicable, in other words, for an "uncapacitated" problem. The program, in this case, will generate the necessary information to conform to formulation (P).

Figure 3 presents a schematic diagram of the deck structure for using ZIPCAP. The cards include both the job control cards and the input data cards required to solve a problem.

Following is a detailed description of these cards (in the same order as presented in Figure 3); it provides the necessary coding information.

(1) Job and JCL Cards: Figures 4A and 4B show two alternatives for the use of JCL cards.



\*When solving a capacitated problem, only the  $s_k$  and  $d_{ijk}$  cards are required; when solving an uncapacitated problem, only the  $e_{ik}$  cards are required.

Figure 3

ZIPCAP Data Structure

**FORTRAN Coding Form**

[illegible][illegible]

**Figure 4A**

### Job and JCL Cards -- Alternative I

IBM

FORTRAN Coding Form

1. NAME		2. PROJECT		3. DATE		4. TIME		5. PAGE	

FORTRAN STATEMENT									
1	2	3	4	5	6	7	8	9	10
//	STANDARD	JOB	CARD						
//	EXEC	F0RG6,	DSN='Y0UR	LIBRARY',	PR0G=Z1PCAP				
//	F0RT.	SYSIN	DD *						
		↑							
	(S0UFC	PF0GRFM)							
		↓							
//									

Step 1: Compile and Store Procedure

FORTRAN STATEMENT									
1	2	3	4	5	6	7	8	9	10
//	STANDARD	JOB	CARD						
//	EXEC	F0RG6,	DSN='Y0UR	LIBRARY',	PR0G=Z1PCAP				
//	G0.	SYSLIB	DD						
//	DD	DSN=GWU.	PLIXLIB,	DISP=SHR					
//	G0.	SYSIN	DD *						
		↑							
	(DATA	CARDS)							
		↓							
//									

} JCL CARDS FOR  
TIMET SUBROUTINE

Step 2: Load and Go Procedure

Figure 4B

Job and JCL Cards -- Alternative II

Alternative I uses the source program for the compilation and execution procedure. In Alternative II, the program is first compiled and stored, and thereafter, for a given problem, only the execution is necessary.

The program uses the TIMET subroutine from the PL1 Library in order to record the elapsed time. This time includes only the execution time of the program and excludes the time needed to read the input and to write the output. The JCL cards needed for this purpose are identified in Figures 4A and 4B. If this subroutine is not available, the corresponding JCL cards are not required. In that case the appropriate source program cards should also be removed. (These cards are identified in the program listing in Appendix A.)

(2) Options Card

The options card is the first input data card and is described in Table 2.

(3) Parameter Card

Table 3 describes the parameter card.

Table 3  
THE PARAMETER CARD

Columns	Format	Name	Definition
1-5	I5	M	Number of designs (m)
6-10	I5	N	Number of activities (n)
11-15	I5	P	Number of facilities(p)



Table 2  
OPTIONS CARD

Columns	Format	Name	Value and Definition	Remarks
1	I1	IINPT	1 if input listing desired 0 otherwise	Input listing includes: - parameters $m, n, p$ - $a_{ij}, b_k, s_k, d_{ijk},$ and $e_{ik}$
2	I1	ICAPR	1 if capacity rule to be used 0 otherwise	When solving a capacitated problem (i.e., where capacity constraints are applicable), set ICAPR=1; for an uncapacitated problem, it should be set to 0.
3	I1	ISTEP	0 if listing of intermediate steps not desired 1 if summary of intermediate steps desired 2 if detailed intermediate steps desired	<ul style="list-style-type: none"> <li>Summary of intermediate steps includes, for each node, node number, lower bound, upper bound (if applicable), best upper bound, and branching variable. (This information is sufficient to construct a branch-and-bound tree.)</li> <li>Detailed intermediate steps list every computational step at every node. This option is useful only when changing/debugging the program.</li> <li>Even with ISTEP = 0, the output includes the total number of nodes explored.</li> </ul>

Continued

Table 2  
(Continued)

Columns	Format	Name	Value and Definition	Remarks
4	I1	IUNCAP	0 if solving a capacitated problem 1 if solving an uncapacitated problem	for IUNCAP=0, ICAPR should be 1, and for IUNCAP=1, ICAPR should be 0. (for IUNCAP=1, if ICAPR=1, the program overrides it and sets ICAPR=0)
5-10	F6.5	EPS	A fractional value if a sub-optimal solution is acceptable. This value should be zero if an optimal solution is desired.	EPS = 0.002 implies that the resulting solution may be suboptimal, but is guaranteed to have a value within 0.2% of the optimal value.
11-20	F10.3	ET	Elapsed time in seconds, if specified, at which the node and bounds information is printed.	<p>Specifying ET is useful in a case where ISTEP = 0 and the program terminates before reaching the final solution.</p> <p>When applicable, the information printed includes:</p> <ul style="list-style-type: none"> <li>- the node being explored at the specified time</li> <li>- the best upper bound, the corresponding solution, and the node at which this solution was found</li> <li>- detailed steps for the node being explored, by assuming ISTEP=2</li> </ul> <p>The above information is useful in assessing the extent of the branch-and-bound tree explored until time ET.</p> <p>This option uses the TIMET subroutine</p>

(4) to (8) Other Input Data Cards

Table 4 describes the cards for the variable costs  $a_{ij}$ , the fixed costs  $b_k$ , the capacity availabilities  $s_k$ , and the capacity usage values  $d_{ijk}$ . However, if solving an uncapacitated problem, i.e., where capacity constraints are not active, the facility-design values  $e_{ik}$  are required as input data instead of the  $s_k$  and  $d_{ijk}$  values.

The input data values for  $a_{ij}$ ,  $b_k$ ,  $s_k$ , and  $d_{ijk}$  are required in integer format. Any real-valued data can be converted into integer values by scaling, and the optimal solution reconverted to real values later. For example, real-valued data of 215.35, 116.50, ..., 180.61 can be treated as 21535, 11650, ..., 18061; and an optimal value of 584230 obtained from the program would correspond to 5842.30.

(9) End Card

This is the usual end card following input data cards. It contains '//' in the first two columns.

5. Test Problems and Output Illustrations5.1 Test Problem 1

5.1.1 Problem formulation (stated for ease of reference)

The problem is to find  $x_{ij}$  and  $y_k$  variables that

Table 4

## OTHER INPUT DATA CARDS

Columns	Format	Name	Definition	Remarks
1-10, 11-20, 21-30, etc.	8I10	A(I,J)	Variable cost of activity j using design i	The order of the input data $a_{ij}$ is $a_{11}, a_{21}, \dots,$ $a_{m1}; a_{12}, a_{22}, \dots,$ $a_{m2}; \dots, a_{1n}, a_{2n}, \dots,$ $a_{mn}$
1-10, 11-20, 21-30, etc.	8I10	B(K)	Fixed cost of facility k	The order for the $b_k$ is $b_1,$ $b_2, \dots, b_p$
1-10, 11-20, 21-30, etc.	8I10	S(K)	Capacity available at facility k	. The order for the $s_k$ is $s_1,$ $s_2, \dots, s_p$ . This input is required when solving a capacitated problem only, i.e., for IUNCAP=0. For an uncapacitated problem, i.e., for IUNCAP=1, this information is generated by the program.

Continued

Table 4  
(Continued)

Columns	Format	Name	Definition	Remarks
1-10, 11-20, 21-30, etc.	8I10	D(I,J,K)	Capacity required at facility $k$ for activity $j$ when activity $j$ uses design $i$ . A $d_{ijk}$ value is positive if design $i$ uses facility $k$ ; otherwise it is 0.	<p>The order of input data is all <math>d_{ijk}</math> for <math>k=1</math>, then all <math>d_{ijk}</math> for <math>k=2</math>, etc., starting on a new card when the value of <math>k</math> changes.</p> <p>The order of input for a facility <math>k</math> is <math>d_{11k}, d_{21k}, \dots, d_{m1k};</math>  <math>d_{12k}, d_{22k}, \dots, d_{m2k};</math>  <math>\dots</math>  <math>d_{1nk}, d_{2nk}, \dots, d_{mnk}.</math></p> <p>This input, similar to <math>s_k</math>, is required for capacitated problems only, i.e., for IUNCAP=0. This information is generated by the program for an uncapacitated problem, i.e., for IUNCAP=1.</p>

Continued

Table 4  
(Continued)

Columns	Format	Name	Definition	Remarks
1,2,3,...	8011	$E(I,K)$	$e_{ik} = 1 \text{ if design } i \text{ uses facility } k$ $= 0 \text{ otherwise}$	<p>The order of input for the <math>e_{ik}</math> is <math>e_{11}, e_{21}, \dots, e_{m1}; e_{12}, e_{22}, \dots, e_{m2}; \dots, e_{1p}, e_{2p}, \dots, e_{mp}</math>.</p> <p>This input is required when solving an uncapacitated problem only, i.e., for IUNCAP=1. The program, then generates <math>d_{ijk}</math> and <math>s_k</math> as follows:</p> $d_{ijk} = 1 \text{ for all } j \text{ if } e_{ik} = 1$ $= 0 \text{ for all } j \text{ if } e_{ik} = 0$ $s_k = n \text{ for all } j$ <p><math>e_{ik}</math> is generated by the program for a capacitated problem, as follows:</p> $e_{ik} = 1 \text{ if } \sum_j d_{ijk} \text{ is positive } (>0)$ $= 0 \text{ otherwise}$

$$\begin{aligned}
 (P) \left\{ \begin{array}{ll} \text{Minimize} & \sum_{j=1}^n \sum_{i=1}^m a_{ij} x_{ij} + \sum_{k=1}^p b_k y_k \quad (1) \\ \text{subject to} & \sum_{i=1}^m x_{ij} = 1 \quad j=1, \dots, n \quad (2) \\ & \sum_{j=1}^n \sum_{i=1}^m d_{ijk} x_{ij} \leq s_k y_k \quad k=1, \dots, p \quad (3) \\ & x_{ij}, y_k = 0 \text{ or } 1, \quad (4) \end{array} \right.
 \end{aligned}$$

where  $m = 3$ ,  $n = 4$ ,  $p = 5$ .

The  $a_{ij}$  values are:

		j			
		1	2	3	4
i	1	9,847	15,718	4,450	8,925
	2	10,192	10,542	4,422	8,118
	3	11,019	9,750	4,609	8,337

$$b_1 = 1750 \quad b_2 = 2000 \quad b_3 = 1750 \quad b_4 = 1350 \quad b_5 = 1000$$

$$s_1 = 400 \quad s_2 = 400 \quad s_3 = 1000 \quad s_4 = 400 \quad s_5 = 400$$

The  $d_{ijk}$  values are as follows:

		j			
k=1		1	2	3	4
i	1	0	0	0	0
	2	186	60	261	148
	3	154	38	175	100

		j			
k=2		1	2	3	4
i	1	0	0	0	0
	2	0	0	0	0
	3	154	38	175	100

		j			
k=3		1	2	3	4
i	1	525	144	1,011	396
	2	186	60	261	148
	3	154	38	175	100

		j			
k=4		1	2	3	4
i	1	0	0	0	0
	2	0	0	0	0
	3	154	38	175	100

		j			
k=5		1	2	3	4
i	1	0	0	0	0
	2	186	60	261	148
	3	154	38	175	100

5.1.2 Coded Input: Figure 5 shows the coded information for test problem 1. Various segments in this figure are numbered in parentheses to correspond to the deck structure of Figure 3.





5.1.3 Annotated Output: The computer output for this test problem is presented in Appendix C. Most segments of this output are self-explanatory. The order of these segments is lettered in circles as follows:

- (a) Options selected -- information given on options card
- (b) m, n, and p -- information from parameter card. This segment is printed if the option IINPT=1
- (c) Other input data, i.e.,  $a_{ij}$ ,  $b_k$ ,  $s_k$  and  $d_{ijk}$  for a capacitated problem; and  $e_{ik}$  as generated by the program. For an uncapacitated problem, input data  $a_{ij}$ ,  $b_k$ , and  $e_{ik}$ ; and  $s_k$  and  $d_{ijk}$  as generated by the program. This segment is printed if the option IINPT=1
- (d) Intermediate steps' display, depending on the value of the option ISTEP. No display for ISTEP=0, summary for ISTEP=1, and detailed steps for ISTEP=2
- (e) Elapsed time in seconds to execute the program, excluding the time to read the input data and to write the input and output. It includes the time to print intermediate steps if option ISTEP=1 or 2
- (f) Total number of nodes explored for the problem solved
- (g) Optimal solution showing the  $x_{ij}$  variables with value 1,  $y_k$  variables with value 1 or 0, and the value of the objective function. If the problem does not have a feasible solution, this segment prints "Problem does not have a feasible solution because the capacity constraints cannot be satisfied."

The node and bounds information at a specified time is displayed only if option RT is assigned a particular value.

## 5.2 Test Problem 2

This is a capacitated problem having 5 designs, 4 activities and 8 facilities.

The problem formulation and coded input are similar to that of test problem 1.

The annotated output (which includes the listing of input data) is presented in Appendix D. The annotations are the same as those already described for test problem 1. Note that the options selected include ISTEP=0, hence no intermediate steps; and EPS=0.002, implying that the solution may be suboptimal within  $\pm 0.2\%$ .

## 5.3 Test Problem 3

Although ZIPCAP is designed for capacitated problems, it may be used to solve uncapacitated problems as a special case. For demonstration, this test problem is uncapacitated, having 10 designs, 8 activities, and 8 facilities. The input data include the options card, the parameter card, input data values for  $a_{ij}$ ,  $b_k$ , and  $e_{ik}$ . The program generates the  $d_{ijk}$  and  $s_k$  values.

Appendix E shows the annotated output for this problem. Note that the options selected include IINPT=1, hence the data listing; ISTEP=0, therefore no intermediate steps; ICAPR=0 and IUNCAP=1 being an uncapacitated problem. Note that option ET has been specified a value of 30.0. Thus, if the optimal solution was not reached within 30 seconds, the node and bounds information at time ET would have been displayed.

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APPENDIX A  
ZIPCAP LISTING

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C          ZIPCAP, A ZERO-ONE INTEGER PROGRAM IS DESIGNED      00000010
C          TO SOLVE MULTIACTIVITY MULTIFACILITY CAPACITY=      00000015
C          CONSTRAINED PROBLEMS HAVING VARIABLE AND FIXED      00000020
C          COSTS. IT ALSO SOLVES UNCAPACITATED PROBLEMS AS A    00000030
C          SPECIAL CASE                                         00000040
0001      INTEGER D(35,35,30), A(35,35), CX(35,35), E(35,30),  00000050
          1      B(30), BSOLX(35), BSOLY(30), FLB(30),FIX(35),FIXI(35), 00000060
          2      FUB(30), S(30), SOLX(35), STX(1225)              00000070
0002      REAL      MINC(35), NMINC(35)                          00000080
0003      DIMENSION C(35,35), DIFBR(35), K72(35), MIND(35)      00000090
0004      INTEGER BR0, BR1, FC, FCUB, P                          00000110
0005      REAL      LOWB, MAXDIF, MINSC                           00000120
C          *****OPTIONS AVAILABLE: IINPT, ICAPP, ISTEP, IUNCAP,EPS 00000130
C          IINPT=1 IF INPUT LISTING DESIRED; 0 OTHERWISE        00000140
C          ICAPP=1 IF CAPACITY RULE TO BE USED; 0 OTHERWISE    00000150
C          ISTEP=0 IF LISTING OF INTERMEDIATE STEPS             00000160
C          NOT DESIRED. ISTEP=1 IF SUMMARY OF BRANCH &          00000170
C          BOUND MODES DESIRED. ISTEP=2 IF DETAILED             00000180
C          LISTING OF INTERMEDIATE STEPS DESIRED.              00000190
C          IUNCAP=1 IF SOLVING AN UNCAPACITATED PROBLEM,        00000200
C          0 OTHERWISE.                                          00000210
C          EPS= A FRACTIONAL VALUE IF SUBOPTIMAL                00000220
C          SOLUTION DESIRED, E.G., EPSILON AS 0.005            00000230
C          IMPLIES SOLUTION TO BE WITHIN +0.5 PERCENT          00000240
C          OF THE OPTIMAL SOLUTION. EPS=0.0 IF OPTIMAL          00000250
C          SOLUTION DESIRED.                                     00000253
C          ET= ELAPSED TIME IN SECONDS, IF SPECIFIED, AT        00000256
C          WHICH THE NODE AND BOUNDS RELATED INFORMATION        00000260
C          IS PRINTED. THIS IS USEFUL IN A SITUATION IF         00000263
C          ISTEP=0 AND THE PROGRAM TERMINATES BEFORE            00000266
C          REACHING THE FINAL SOLUTION.                          00000270
C          *****READ INPUT DATA*****                         00000273
0006      READ 10, IINPT, ICAPP, ISTEP, IUNCAP, EPS, ET          00000280
0007      10 FORMAT (4I1, F6.5, F10.3)                          00000290
C          M= NUMBER OF DESIGNS                                  00000300
C          N= NUMBER OF ACTIVITIES                              00000310
C          P= NUMBER OF FACILITIES                              00000320
0008      RE-D 20,M,N,P                                          00000330
0009      20 FORMAT (3I5)                                         00000340
C          A(I,J): VARIABLE COST MATRIX                        00000350
0010      READ 30, ((A(I,J), I=1,M),J=1,N)                      00000360
0011      30 FORMAT (8I10)                                         00000370
C          B(K): FIXED COST VECTOR                             00000380
0012      READ 30, (B(K),K=1,P)                                    00000390
0013      IF (IUNCAP.EQ.1) GO TO 40                                00000400
C          S(K): CAPACITY LIMIT VECTOR; REQUIRED ONLY          00000410
C          IF IUNCAP=0                                          00000420
0014      READ 30, (S(K),K=1,P)                                    00000430
C          D(I,J,K): CAPACITY USAGE MATRIX; REQUIRED           00000440
C          ONLY IF IUNCAP=0                                     00000450
0015      DO 32 K=1,P                                             00000460
0016      READ 30,((D(I,J,K), I=1,M),J=1,N)                     00000470
0017      32 CONTINUE                                             00000480
0018      DO 37 K=1,P                                             00000490
0019      DO 37 I=1,M                                             00000500
0020      IF (D(I,I,K).EQ.0) GO TO 35                             00000510
0021      E(I,K)=1                                               00000520
0022      GO TO 37                                                 00000530

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0023      35      E(I,K)=0      00000540
0024      37 CONTINUE      00000550
0025      GO TO 90      00000560
      C      E(I,K): DESIGN-FACILITY MATRIX; REQUIRED ONLY      00000570
      C      IF IUNCAP=1      00000580
0026      40 READ 45,((E(I,K),I=1,M),K=1,P)      00000590
0027      45 FORMAT (80I1)      00000600
0028      DO 80 K=1,P      00000610
0029      S(K)=N      00000620
0030      DO 75 I=1,M      00000630
0031      IF (E(I,K).EQ.1) GO TO 65      00000640
0032      DO 60 J=1,N      00000650
0033      D(I,J,K)=0      00000660
0034      60 CONTINUE      00000670
0035      GO TO 75      00000680
0036      65 DO 70 J=1,N      00000690
0037      D(I,J,K)=1      00000700
0038      70 CONTINUE      00000710
0039      75 CONTINUE      00000720
0040      80 CONTINUE      00000730
      C      *****PRINT INPUT DATA*****      00000740
0041      90 PRINT 95, IINPT, ICAPR, ISTEP, IUNCAP, EPS, ET      00000750
0042      95 FORMAT ( '1', ' OPTIONS SELECTED : IINPT=', I1,      00000760
      1      ' ICAPR=', I1, ' ISTEP=', I1, ' IUNCAP=', I1,      00000770
      2      ' EPS=', F8.5, ' ET=', F10.3//')      00000780
0043      IF (IINPT.EQ.0) GO TO 168      00000790
0044      PRINT 100,M,N,P      00000800
0045      100 FORMAT ( '0', T55, 'INPUT DATA', /1X, T55, '-----', /1X,      00000810
      1T41, 'NUMBER OF DESIGNS (M)=', 4X, I4//1X, T41,      00000820
      2'NUMBER OF ACTIVITIES (N)=', 1X, I4//1X, T41,      00000830
      3'NUMBER OF FACILITIES (P)=', 1X, I4//')      00000840
0046      PRINT 105      00000850
0047      105 FORMAT ( 4X, 'VARIABLE COST MATRIX A(I,J)', /4X,      00000860
      1'-----', /)      00000870
0048      DO 110 I=1,M      00000880
0049      110 PRINT 115, I, (A(I,J),J=1,N)      00000890
0050      115 FORMAT ( '0', T6, 'I=', I3, 4X, 8I13, 4(/, 14X, 8I13))      00000900
0051      PRINT 120      00000910
0052      120 FORMAT( '0', //4X, 'FIXED COST VECTOR B(K)', /4X,      00000920
      1'-----', /)      00000930
0053      PRINT 122, (B(K),K=1,P)      00000940
0054      122 FORMAT ( '0', T15, 8I13, 3(/, 14X, 8I13))      00000950
0055      PRINT 125      00000960
0056      125 FORMAT( '0', //4X, 'CAPACITY LIMIT VECTOR S(K)', /4X,      00000970
      1'-----', /)      00000980
0057      PRINT 128, (S(K),K=1,P)      00000990
0058      128 FORMAT ( '0', T15, 8I13, 3(/, 14X, 8I13))      00010000
0059      PRINT 130      00010100
0060      130 FORMAT( '0', //4X, 'CAPACITY USAGE MATRIX D(I,J,K)', /4X,      00010200
      1'-----', /)      00010300
0061      DO 150 K=1,P      00010400
0062      PRINT 135,K      00010500
0063      135 FORMAT ( '0', //5X, 'K=', I3//)      00010600
0064      DO 145 I=1,M      00010700
0065      PRINT 140,I,(D(I,J,K), J=1,N)      00010800
0066      140 FORMAT ( '0', T6, 'I=', I3, 4X, 8I13, 4(/, 14X, 8I13))      00010900
0067      145 CONTINUE      00011000
0068      150 CONTINUE      00011100

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0069      PRINT 155
0070      155 FORMAT('0',//4X,'DESIGN=FACILITY MATRIX E(I,K)',//4X,
1          '-----',//)
0071      DO 160 I=1,M
0072      PRINT 158, I, (E(I,K),K=1,P)
0073      158 FORMAT('0', T6, 'I=', I3, 4X, 8I13, 3(/, 14X, 8I13))
0074      160 CONTINUE
0075      168 IF (ISTEP.EQ.0) GO TO 190
0076      IF (ISTEP.EQ.1) GO TO 175
0077      PRINT 170
0078      170 FORMAT('0',///55X,'DETAILED LISTING OF STEPS',/)
0079      GO TO 190
0080      175 PRINT 180
0081      180 FORMAT('0',///55X,'SUMMARY OF STEPS',/)
0082      C *****INITIALIZE*****
0083      190 BUB=9999999.
0084      BUB$= BUB/ (1.0+EPS)
0085      NSX=0
0086      MOD=1
0087      IBMOD=1
0088      IN-T=0
0089      INSET=0
0090      DO 205 J=1,M
0091      FIX(J)=0
0092      KT2(J)=0
0093      DO 205 I=1,M
0094      CX(I,J)=0
0095      205 CONTINUE
0096      LQ1=0
0097      LQ2=0
0098      LR2=0
0099      CALL TIMET(1TO)
0099      IF(NSX.EQ.0) GO TO 283
C          CX(I,J) CONTAINS FIXED AND FREE X(I,J) VARIABLES.
C          STX(INS) CONTAINS FIXED X(I,J) VARIABLES.
C          CX(I,J) AND STX(INS) ARE UPDATED BY THE CAPACITY
C          RULE, THE BOUNDING RULE, AND THE RULE FOR
C          BRANCHING AND BACKTRACKING.
C          IN CX(I,J) A FIXED VARIABLE IS RECORDED AS 1 OR
C          2, AND A FREE VARIABLE AS 0.
C          A VALUE OF 1 IMPLIES THAT THAT PARTICULAR VARIABLE
C          IS FIXED, AND FIX(J) IS SET EQUAL TO 1 IMPLYING
C          THAT COLUMN J HAS A FIXED VARIABLE OF VALUE 1.
C          A VALUE OF 2 IMPLIES THAT THAT PARTICULAR VARIABLE
C          SHOULD NOT BE CONSIDERED FOR CURRENT COMPUTATIONS.
C          AN X(I,J) RECORDED IN CX(I,J) AS 1 DUE TO THE
C          BRANCHING RULE IS RECORDED IN STX(INS) AS X*100+J.
C          AN X(I,J) RECORDED IN CX(I,J) AS 1 DUE TO THE
C          CAPACITY RULE OR THE BOUNDING RULE IS RECORDED IN
C          STX(INS) AS (X*100+J)+1000000.
C          AN X(I,J) RECORDED IN CX(I,J) AS 2 IS RECORDED IN
C          STX(INS) AS -(X*100+J)-1000000.
0100      210 IF (ISTEP.EQ.0) GO TO 225
0101      215 PRINT 220,MOD
0102      220 FORMAT('0',// 6X,'NODE NUMBER',I5/)
C          *****UPDATE CX(I,J) FOR BRO*****
C          BRO IS THE RIGHT BRANCHING VARIABLE
0103      225 LX=BRO

```

\*Relevant to the TIMET Subroutine



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0104      IX=LX/100
0105      JX=LX-IX*100
0106      CX(IX,JX)=2
0107      KT2(JX)=KT2(JX)+1
0108      FIX(JX)=0
0109      LQ1=LQ1+1
0110      IF (KT2(JX).LT.(M-1)) GO TO 270
0111      DO 255 I=1,M
0112          IF (CX(I,JX).EQ.2) GO TO 255
0113          CX(I,JX)=1
0114          NSX=NSX+1
0115          STX(NSX)= (I*100+JX)+1000000
0116          FIX(JX)=1
0117          LQ1=LQ1+1
0118          FIXI(JX)=I
0119          GO TO 270
0120      255 CONTINUE
0121      270 LQ2=0
0122          LR2=0
0123          GO TO 283
0124      272 IF (ISTEP.EQ.0) GO TO 276
0125      PRINT 220,NOD
C          *****UPDATE CX(I,J) FOR BRI*****
C          BRI IS THE LEFT BRANCHING VARIABLE
0126      276 LQ2=0
0127          LR2=0
0128          LX=BRI
0129          IX=LX/100
0130          JX=LX-IX*100
0131          CX(IX,JX)=1
0132          FIX(JX)=1
0133          LQ1=LQ1+1
0134          DO 279 I=1,M
0135              IF (IX.EQ.I) GO TO 281
0136      279 CONTINUE
0137      281 FIXI(JX)=IX
0138      283 IF (ISTEP.NE.2) GO TO 300
0139      285 DO 295 I=1,M
0140          PRINT 290, I,(CX(I,J),J=1,N)
0141      290 FORMAT (/5X,'CX(I,J)',4X,'I=',I3,2X, 2014/23X, 2014)
0142      295 CONTINUE
0143          PRINT 297,(FIX(I),J=1,N)
0144      297 FORMAT (/5X,'FIX(I)',12X, 2014/23X, 2014)
C          *****APPLY CAPACITY RULE*****
C          AND UPDATE CX(I,J) AND STX(INS).
0145      300 IF (IUNCAP.EQ.1) GO TO 2300
0146      305 IF (ICAPR.EQ.0) GO TO 2300
0147      310 DO 2000 K=1,P
C          FIND THE SUM OF MINIMUM D(I,J,K) OVER EACH J FOR A
C          GIVEN K, I.E., MINSO= SUM OF MIND(J)
0148      MINSO=0
0149      DO 900 J=1,N
0150          IF(FIX(J).EQ.0) GO TO 350
C          IF FIX(J)=1, SET MIND(J)=D(I,J,K) FOR CX(I,J)=1
C          AND MOVE TO NEXT COLUMN J
0151      IND(=FIXI(J)
0152      MIND(J)=D(IND1,J,K)
0153      GO TO 800

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0154      350      LK=0                                00002160
0155      I=1                                           00002170
0156      MIND(J)=D(I,J,K)                             00002180
C          SKIP D(I,J,K) WHEN CX(I,J)=2 & MOVE TO NEXT ROW I 00002190
0157      400      IF(CX(I,J).EQ.2) GO TO 600          00002200
0158      500      IF(D(I,J,K).LT.MIND(J)) MIND(J)=D(I,J,K) 00002210
0159      GO TO 700                                     00002220
0160      600      LK=LK+1                             00002230
0161      IF(I.GT.LK) GO TO 700                       00002240
0162      I=I+1                                         00002250
0163      MIND(J)=D(I,J,K)                             00002260
0164      GO TO 750                                     00002270
0165      700      I=I+1                                00002280
0166      750      IF(I.LE.M) GO TO 400                 00002290
0167      800      MINSO=MINSO+MIND(J)                 00002300
0168      900      CONTINUE                             00002310
0169      910      IF (ISTEP.NE.2) GO TO 980            00002320
0170      PRINT 950, K, MINSO, (MIND(J),J=1,N)         00002330
0171      950      FORMAT('0',*K, MINSO, (MIND(J),J=1,N)*,10I10) 00002340
C          FIND BALANCE AVAILABLE CAPACITY IBALD FOR A GIVEN K 00002350
C          IF IBALD IS NEGATIVE, THEN BACKTRACK.         00002360
0172      980      IBALD=S(K)-MINSO                    00002380
0173      1000     IF (IBALD.LT.0) GO TO 6200           00002390
0174      DO 1500 J=1,N                                00002400
C          SKIP COLUMN J IF FIX(J)=1                    00002410
0175      IF (FIX(J).EQ.1) GO TO 1500                  00002420
0176      DO 1300 I=1,M                                00002430
C          SKIP ROW I IF CX(I,J)=2                     00002440
0177      1100     IF(CX(I,J).EQ.2) GO TO 1300          00002450
C          COMPUTE DIFFERENCE BETWEEN D(I,J,K) AND MIND(J). 00002470
C          IF IT IS MORE THAN AVAILABE BALANCE, SET CX(I,J)=2 00002480
0178      1200     IDIFD=D(I,J,K)-MIND(J)              00002490
0179      IF ((IDIFD-IBALD).LE.0) GO TO 1300            00002510
0180      CX(I,J)=2                                     00002520
0181      NSX=NSX+1                                     00002523
0182      STX(NSX)=-(I*100+J)-1000000                 00002526
C          LQ2 COUNTS THE NUMBER OF CX(I,J) VALUES SET EQUAL 00002530
C          TO 2 IN A CYCLE                              00002540
0183      LQ2=LQ2+1                                    00002550
C          KT2(J) KEEPS AN ACCOUNT OF CX(I,J) VALUES SET EQUAL 00002560
C          TO 2 FOR COLUMN J                            00002570
0184      KT2(J)=KT2(J)+1                              00002580
C          FOR COLUMN J, IF ALL BUT ONE CX(I,J) VALUES ARE 00002590
C          EQUAL TO 2, SET THAT CX(I,J)=1 & SET FIX(J)=1 00002600
0185      IF(KT2(J).LT.(M-1)) GO TO 1300              00002610
0186      DO 1250 LR=1,M                                00002620
0187      IF(CX(LR,J).EQ.2) GO TO 1250                 00002630
0188      CX(LR,J)=1                                    00002640
0189      NSX=NSX+1                                     00002643
0190      STX(NSX)= (LR*100+J)+1000000                 00002646
0191      FIX(J)=1                                      00002650
C          LQ1 KEEPS AN ACCOUNT OF COLUMNS FOR WHICH FIX(J)=1 00002655
0192      LQ1=LQ1+1                                    00002660
C          FIXI(J) SPECIFIES INDEX I FOR WHICH FIX(J)=1    00002662
0193      FIXI(J)=LR                                   00002665
0194      GO TO 1500                                    00002670
0195      1250     CONTINUE                             00002680
0196      1300     CONTINUE                             00002690

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0197      1500 CONTINUE                                00002700
0198      1800 IF (ISTEP.NE.2) GO TO 2000              00002710
0199      PRINT 1900, K, LQ2, LQ1                      00002720
0200      1900 FORMAT ('0', 'K=', I3, ' LQ2=', I3, ' LQ1=', I3) 00002730
0201      DO 1930 I=1, M                                00002740
0202          PRINT 290, I, (CX(I, J), J=1, N)          00002750
0203      1930 CONTINUE                                00002770
0204      PRINT 297, (FIX(J), J=1, N)                  00002780
0205      2000 CONTINUE                                00002800
C          A CYCLE EXAMINES ALL THE FACILITIES.        00002803
C          IF IN A CYCLE, THE CAPACITY RULE RESULTS IN SETTING 00002810
C          ADDITIONAL CX(I, J) VALUES EQUAL TO 2, THEN REPEAT 00002820
C          THE CYCLE. BUT IF FIX(J)=1 FOR ALL J, THEN DO NOT 00002830
C          REPEAT THE CYCLE.                            00002835
0206      IF (LQ1.EQ.N) GO TO 2300                      00002840
0207      IF (LQ2.EQ.LR2) GO TO 2300                    00002845
0208      2200 LR2=LQ2                                  00002860
0209      GO TO 300                                      00002870
C          *****SOLVE (LAGRANGIAN) RELAXED PROBLEM***** 00002880
C          FIND FLB(K) — VECTOR OF FACILITIES FOR COMPUTING 00002890
C          C(I, J) MATRIX & LOWER BOUND. IT HAS VALUE 1 IF A 00002900
C          FACILITY IS USED, OTHERWISE IT HAS 0 VALUE.    00002910
0210      2300 DO 2500 K=1, P                            00002920
0211          FLB(K)=0                                    00002930
0212      2500 CONTINUE                                00002940
0213          DO 3000 J=1, N                            00002950
0214              IF (FIX(J).EQ.0) GO TO 3000            00002960
0215              INDI=FIXI(J)                          00002970
0216          DO 2550 K=1, P                            00002990
0217              IF (E(INDI, K).EQ.0) GO TO 2550        00003000
0218              IF (FLB(K).EQ.1) GO TO 2550            00003010
0219              FLB(K)=1                               00003020
0220      2550 CONTINUE                                00003030
0221      3000 CONTINUE                                00003060
0222          IF (ISTEP.NE.2) GO TO 3150                00003070
0223          PRINT 3100, (FLB(K), K=1, P)              00003080
0224          3100 FORMAT('0', '(FLB(K), K=1, P) ', 20I4/16X, 20I4) 00003090
C          COMPUTE COST MATRIX C(I, J) FOR THE RELAXED PROBLEM 00003100
0225      3150 DO 3400 J=1, N                            00003110
0226          DO 3300 I=1, M                            00003120
0227              BSUM=0.0                               00003130
0228          DO 3200 K=1, P                            00003140
0229              IF (FLB(K).EQ.1) GO TO 3200            00003150
0230              IF (E(I, K).EQ.0) GO TO 3200          00003160
0231              BSUM=BSUM+(B(K) * (FLOAT(D(I, J, K))/ FLOAT(S(K)))) 00003170
0232      3200 CONTINUE                                00003180
0233          C(I, J)=A(I, J)+BSUM                      00003190
0234      3300 CONTINUE                                00003200
0235      3400 CONTINUE                                00003210
0236          IF (ISTEP.NE.2) GO TO 3445                00003220
0237          DO 3430 I=1, M                            00003230
0238              PRINT 3420, I, (C(I, J), J=1, N)        00003250
0239          3420 FORMAT (1/5X, 'C(I, J)', 5X, 'I=', I3, 2X, 5F15.4, 00003260
1              6(1/23X, 5F15.4))                      00003265
0240      3430 CONTINUE                                00003270
C          FIND SUM OF MINIMUM C(I, J) VALUES OVER EACH J, 00003290
C          I.E., MINSC=SUM OF MINC(J).                  00003300
C          IF FIX(J)=1, THEN MINC(J)=C(I, J) WHERE CX(I, J)=1 00003310

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0241      3445 MINSC=0.0                                00003320
0242      DO 3900 J=1,N                                00003340
0243          IF (FIX(J).EQ.0) GO TO 3500                00003350
0244          INDI=FIXI(J)                                00003360
0245          MINC(J)=C(INDI,J)                          00003370
0246          SOLX(J)=INDI                                00003380
0247          GO TO 3850                                  00003410
0248      3500      LK=0                                  00003430
0249          I=1                                          00003440
C          SKIP C(I,J) ELEMENT IF CX(I,J)=2 & MOVE TO NEXT I 00003470
0250      3550      IF (CX(I,J).EQ.2) GO TO 3700          00003480
0251          IF ((I=LK).EQ.1) GO TO 3600                00003485
0252          IF (C(I,J).GE.MINC(J)) GO TO 3750          00003490
0253      3600      MINC(J)=C(I,J)                        00003500
0254          IMIN=I                                       00003510
0255          GO TO 3750                                  00003520
0256      3700      LK=LK+1                                00003530
0257      3750      I=I+1                                  00003590
0258      3800      IF (I.LE.M) GO TO 3550                00003600
0259          SOLX(J)=IMIN                                00003610
0260      3850      MINSC=MINSC+MINC(J)                  00003620
0261      3900      CONTINUE                              00003630
0262          IF (ISTEP.NE.2) GO TO 3940                  00003640
0263          DO 3920 J=1,N                                00003650
0264              PRINT 3910, J,MINC(J),SOLX(J)           00003660
0265      3910      FORMAT ('0', 'J,MINC(J),SOLX(J)', I5,F15.4,I6) 00003670
0266      3920      CONTINUE                              00003680
C          COMPUTE FIXED COST FC FOR LOWER BOUND          00003710
0267      3940      FC=0                                    00003720
0268          DO 4000 K=1,P                                00003730
0269              IF (FLB(K).EQ.0) GO TO 4000              00003740
0270              FC=FC+B(K)                                00003750
0271      4000      CONTINUE                              00003760
C          *****COMPUTE LOWER BOUND LOWB*****          00003770
0272      4050      LOWB=MINSC+FC                          00003780
0273          IF (ISTEP.EQ.0) GO TO 4150                  00003790
0274          PRINT 4120, MINSC, FC, LOWB                 00003800
0275      4120      FORMAT ('0', ' MINSC, FC, LOWB ', F15.4, I15, F15.4) 00003810
C          COMPARE LOWER BOUND WITH BEST UPPER BOUND STAR 00003820
C          BUBS WHICH EQUALS BUB/(1+EPS). IF LOWB IS     00003830
C          GREATER THAN OR EQUAL TO BUBS, THEN BACKTRACK 00003840
0276      4150      IF (LOWB.GE.BUBS) GO TO 6200           00003850
C          CHECK IF CURRENT SOLUTION SATISFIES CAPACITY  00003880
C          CONSTRAINTS                                    00003890
0277      4200      IF (IUNCAP.EQ.1) GO TO 4420           00003900
0278      4210      DO 4400 K=1,P                          00003910
0279          NSUMD=0                                       00003920
0280          DO 4300 J=1,N                                00003930
0281              IX=SOLX(J)                                00003950
0282              NSUMD=NSUMD+D(IX,J, K)                  00003960
0283      4300      CONTINUE                              00003970
0284          IF (ISTEP.NE.2) GO TO 4320                  00003980
0285          PRINT 4310, K,NSUMD                          00003990
0286      4310      FORMAT ('0', 'K,NSUMD', 2I10)          00004000
0287      4320      IF(NSUMD.LE.S(K)) GO TO 4400           00004010
0288          GO TO 5100                                   00004020
0289      4400      CONTINUE                              00004030
C          *****COMPUTE UPPER BOUND UPB IF CAPACITY CONSTRAINTS 00004040

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C          ARE SATISFIED.                                00004050
C          UPB=SUM OF A(I,J)+FIXED COST FCUB BASED ON      00004060
C          SOLUTION VECTOR SOLX(J)                        00004070
C          VECTOR OF FACILITIES FOR UPPER BOUND FUB(K) HAS 00004080
C          VALUES 1 OR 0 BASED ON FACILITY USED OR OTHERWISE 00004090
0290      4420 DO 4450 K=1,P                                00004100
0291          FUB(K)=0                                      00004110
0292      4450 CONTINUE                                     00004120
0293          NSUMA=0                                       00004130
0294          FCUB=0                                        00004140
0295      4500 DO 4650 J=1,N                                00004150
0296          IX=SOLX(J)                                    00004170
0297          NSUMA=NSUMA+A(IX,J)                          00004180
0298      4550 DO 4600 K=1,P                                00004190
0299          IF (E(IX,K).EQ.0) GO TO 4600                 00004200
0300          IF (FUB(K).EQ.1) GO TO 4600                 00004210
0301          FUB(K)=1                                      00004220
0302          FCUB=FCUB+B(K)                                00004230
0303      4600 CONTINUE                                     00004240
0304      4650 CONTINUE                                     00004250
0305          IF (ISTEP.NE.2) GO TO 4700                   00004260
0306          PRINT 4660, (FUB(K),K=1,P)                  00004270
0307      4660 FORMAT('0',*(FUB(K),K=1,P) ' ', 20I4/16X,20I4) 00004280
0308      4700 UPB=NSUMA+FCUB                                00004290
0309      4708 IF (ISTEP.EQ.0) GO TO 4750                  00004300
0310          PRINT 4710, NSUMA, FCUB, UPB, BUB, BUBS      00004310
0311      4710 FORMAT('0',*NSUMA, FCUB, UPB, BUB, BUBS ' ',2I10, 3F15.4) 00004320
C          COMPARE UPPER BOUND WITH BEST UPPER BOUND      00004330
C          IF UPB IS LESS THAN BUB, SET IT AS BUB AND      00004340
C          NOTE THE SOLUTION                                00004350
0312      4750 IF (UPB.GE.BUB) GO TO 5100                  00004360
0313      4770 BUB=UPB                                       00004370
0314          BUBS= BUB/ (1.0+EPS)                         00004380
0315          IBNOD=NOD                                     00004385
0316          DO 4800 J=1,N                                00004390
0317      4800 BSOLX(J)=SOLX(J)                             00004400
0318          DO 4850 K=1,P                                00004410
0319      4850 BSOLY(K)=FUB(K)                             00004420
C          *****COMPARE LOWB WITH BUBS. IF LOWB IS GREATER 00004430
C          THAN OR EQUAL TO BUBS, THEN BACKTRACK           00004440
0320      4900 IF (LOWB.GE.BUBS)GO TO 6200                 00004450
C          *****IF FIX(J) VALUES ARE 1 FOR EACH J, THEN BACKTRACK 00004480
0321      5100 IF (LQ1.EQ.N) GO TO 6200                   00004500
C          *****APPLY THE BOUNDING RULE*****            00004510
C          IF THE DIFFERENCE BETWEEN C(I,J) AND MINC(J) IS 00004515
C          GREATER THAN THE DIFFERENCE BETWEEN BUBS AND   00004520
C          LOWB, THEN CX(I,J)=2                            00004525
C          *****APPLY BRANCHING RULE AND FIND BRI, THE NEXT 00004530
C          VARIABLE FOR LEFT BRANCHING.                   00004540
C          FIND NMINC(J), THE NEXT HIGHER VALUE THAN MINC(J) 00004550
C          AND DIFBR(J), THE DIFFERENCE BETWEEN THEM.      00004555
0322          OBOUND=BUBS-LOWB                             00004568
0323      5200 DO 5250 J=1,N                                00004570
0324          NMINC(J)=0.0                                  00004580
0325          DIFBR(J)=0.0                                  00004590
0326      5250 CONTINUE                                     00004600
0327          DO 5600 J=1,N                                00004610
C          SKIP TO NEXT J IF FIX(J)=1                    00004620

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0328          IF (FIX(J).EQ.1) GO TO 5600          00004630
0329          LK=0                                00004640
0330          I=1                                  00004650
C              SKIP C(I,J) IF CX(I,J)=2 & MOVE TO NEXT I 00004670
0331      5300  IF (CX(I,J).EQ.2) GO TO 5350          00004680
0332          IF (I.EQ.SOLX(J)) GO TO 5350          00004690
0333          IF ((C(I,J)-MINC(J)).GT.DBOUND) GO TO 5330 00004700
0334          IF ((I-LK).EQ.1) GO TO 5320          00004710
0335          IF (C(I,J).GE.NMINC(J)) GO TO 5400      00004720
0336      5320  NMINC(J)=C(I,J)                    00004730
0337          GO TO 5400                            00004735
0338      5330  CX(I,J)=2                          00004740
0339          NSX=NSX+1                             00004742
0340          STX(NSX)=-(I*100+J)-1000000          00004745
0341          KT2(J)=KT2(J)+1                      00004747
0342          IF (KT2(J).LT.(M-1)) GO TO 5350      00004750
0343          INDI=SOLX(J)                         00004752
0344          CX(INDI,J)=1                        00004755
0345          NSX=NSX+1                             00004758
0346          STX(NSX)= (INDI*100+J)+1000000      00004760
0347          FIX(J)=1                             00004762
0348          LQ1=LQ1+1                             00004764
0349          FIXI(J)=INDI                        00004766
0350          GO TO 5600                           00004768
0351      5350  LK=LK+1                             00004770
0352      5400  I=I+1                               00004775
0353          IF (I.LE.M) GO TO 5300              00004780
0354      5500  DIFBR(J)=NMINC(J)-MINC(J)          00004785
0355      5600  CONTINUE                           00004790
0356          IF (ISTEP.NE.2) GO TO 5650          00004795
0357          DO 5620 I=1,M                      00004820
0358              PRINT 290, I, (CX(I,J),J=1,N)    00004830
0359      5620  CONTINUE                           00004850
0360          PRINT 297, (FIX(J),J=1,N)            00004860
C              IF FIX(J)=1 FOR ALL J, THEN BACKTRACK. 00004880
0361      5650  IF (LQ1.EQ.N) GO TO 6200          00004890
C              FIND MAXDIF, THE MAXIMUM DIFFERENCE DIFBR(J) 00004900
0362          LF=0                                00004905
0363          DO 5800 J=1,N                       00004910
0364              IF (FIX(J).EQ.1) GO TO 5690      00004915
0365              IF ((J-LF).EQ.1) GO TO 5660      00004920
0366              IF (DIFBR(J).LT.MAXDIF) GO TO 5800 00004925
0367      5660  MAXDIF=DIFBR(J)                   00004930
0368              LJ=J                             00004935
0369              GO TO 5800                       00004940
0370      5690  LF=LF+1                           00004943
0371      5800  CONTINUE                           00004946
0372          IF (ISTEP.NE.2) GO TO 5840          00004950
0373          DO 5820 J=1,N                       00004953
0374              IF (FIX(J).EQ.1) GO TO 5820      00004956
0375              PRINT 5810, J, NMINC(J), MINC(J), DIFBR(J) 00004960
0376      5810  FORMAT ('0','J,NMINC(J),MINC(J),DIFBR(J)', I5,3F15.4) 00004963
0377      5820  CONTINUE                           00004966
C          *****BRANCHING VARIABLE BR1 CORRESPONDS TO MAXDIF***** 00004970
0378      5840  DO 5900 J=1,N                     00004980
0379              IF (J.NE.LJ) GO TO 5900          00004990
0380      5850  BR1=SOLX(J)*100+J                 00005000
0381          IF (ISTEP.EQ.0) GO TO 6020          00005010

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0382          PRINT 5880, BR1
0383          5880  FORMAT('0', ' BR1', I10)
0384          GO TO 6020
0385          5900 CONTINUE
C          *****UPDATE STX(INS) AND NSX*****
C          NSX REPRESENTS THE NUMBER OF VARIABLES IN STX(INS)
0386          6020 NSX=NSX+1
0387          6040 STX(NSX)=BR1
0388          IF (ISTEP.NE.2) GO TO 6100
0389          PRINT 6088, (STX(INS), INS=1, NSX)
0390          6088  FORMAT('0', ' STX(INS)', 10I10, 122(/, 10X, 10I10))
C          *****MOVE TO THE NEXT (LEFT BRANCH) NODE AND APPLY
C          CAPACITY RULE
0391          6100 NOD=NOD+1
0392          6110 IF (ET.EQ.0.0) GO TO 6150
0393          IF (INSET.EQ.1) GO TO 6147
0394          IF (INET.EQ.1) GO TO 6150
0395          CALL TIMET(INT)
0396          ELTN=(INT-ITO)*26.04E-6
0397          IF (ELTN.LT.ET) GO TO 6150
0398          6120 PRINT 6125, NOD, ELTN, BUB, BUBS, IBMOD
0399          6125  FORMAT('0', ' WAS AT NODE', I6, ' AT ELAPSED TIME =', F10.4,
1              ' SECONDS.', /1X, ' BUB=', F15.4, ' BUBS=', F15.4,
2              ' AT NODE=', I7)
0400          IBUB=BUB
0401          IF (IBUB.EQ.9999999) GO TO 6146
0402          6130 PRINT 6135, (BSOLX(J), J=1, N)
0403          6135  FORMAT('0', ' SOLUTION CORRESPONDING TO BUB IS', /1X,
1              ' (BSOLX(J), J=1, N)', 10I8, 3(/10X, 10I8))
0404          6140 PRINT 6145, (BSOLY(K), K=1, P)
0405          6145  FORMAT(/1X, ' (BSOLY(K), K=1, P)', 10I8, 2(/10X, 10I8))
0406          6146 INET=1
0407          INIS=ISTEP
0408          INSET=1
0409          ISTEP=2
0410          GO TO 6150
0411          6147 ISTEP=INIS
0412          INSET=0
0413          6150 GO TO 272
C          *****END IF AT THE ROOT NODE*****
0414          6200 IF (NSX.EQ.0) GO TO 8100
0415          6220 IF (IABS(STX(NSX)).GT.1000000) GO TO 6500
0416          6250 BRO=STX(NSX)
0417          6270 STX(NSX)=BRO-1000000
0418          IF (ISTEP.EQ.0) GO TO 6308
0419          PRINT 6305, BRO
0420          6305  FORMAT('0', ' BRO ', I10)
0421          6308 IF (ISTEP.NE.2) GO TO 6330
0422          PRINT 6088, (STX(INS), INS=1, NSX)
C          *****MOVE TO THE NEXT (RIGHT BRANCH) NODE AND APPLY
C          CAPACITY RULE
0423          6330 NOD=NOD+1
0424          6410 IF (ET.EQ.0.0) GO TO 6450
0425          IF (INSET.EQ.1) GO TO 6445
0426          IF (INET.EQ.1) GO TO 6450
0427          CALL TIMET(INT)
0428          ELTN=(INT-ITO)*26.04E-6
0429          IF (ELTN.LT.ET) GO TO 6450

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\*Relevant to the TIMET Subroutine

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0430      6420 PRINT 6125, MOD, ELTN, BUB, BUBS, IBMOD      00005528*
0431      IBUB=BUB      00005530*
0432      IF (IBUB.EQ.9999999) GO TO 6442      00005532*
0433      6430 PRINT 6135, (BSOLX(J),J=1,N)      00005533*
0434      6440 PRINT 6145, (BSOLY(K), K=1,P)      00005536*
0435      6442 INET=1      00005538*
0436      INIS=ISTEP      00005540*
0437      INSET=1      00005542*
0438      ISTEP=2      00005544*
0439      GO TO 6450      00005546*
0440      6445 ISTEP=INIS      00005548*
0441      INSET=0      00005550*
0442      6450 GO TO 210      00005552
0443      6500 IF (STX(NSX).GT.1000000) GO TO 6520      00005555
0444      LX=STX(NSX)-1000000      00005560
0445      IX=LX/100      00005570
0446      JX=LX-IX*100      00005580
0447      CX:IX,JX)=0      00005590
0448      KT2(JX)=KT2(JX)-1      00005595
0449      GO TO 6550      00005600
0450      6520 LX= STX(NSX)-1000000      00005610
0451      IX=LX/100      00005620
0452      JX=LX-IX*100      00005630
0453      CX(IX,JX)=0      00005640
0454      FIX(JX)=0      00005650
0455      LQ1=LQ1-1      00005660
0456      6550 NSX=NSX-1      00005660
0457      GO TO 6200      00005700
C      *****PRINT THE OUTPUT*****      00005730
0458      8100 IBUB=BUB      00005740
0459      CALL TIMET(IT1)      00005750*
0460      ELT1=(IT1-IT0)*26.04E-6      00005760*
0461      PRINT 8105, ELT1      00005770*
0462      8105 FORMAT ('0',///1X, 'ELAPSED TIME IN SECONDS=', F15.8)      00005780*
0463      PRINT 8120, MOD      00005790
0464      8120 FORMAT ('0', 'TOTAL NUMBER OF NODES EXPLORED =', I3)      00005800
0465      IF (IBUB.EQ.9999999) GO TO 8350      00005810
0466      8130 PRINT 8150      00005820
0467      8150 FORMAT ('0', 'NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN',
1          ' I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.',
2          '/7X, '2. IF EPSILON EPS WAS ASSIGNED A POSITIVE',
3          ' (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.',/)      00005830
0468      8180 PRINT 8200, (BSOLX(J),J=1,N)      00005840
0469      8200 FORMAT('0',T55, 'OPTIMAL SOLUTION',/1X,T55,
1          '-----',/1X, 'X(I,J) WITH VALUE 1:',10I8,
2          '3(/,21X,10I8))      00005850
0470      8220 PRINT 8250, (BSOLY(K), K=1,P)      00005860
0471      8250 FORMAT ('0', 'Y(K) VALUES:', 8X, 10I8, 2(/,21X,10I8))      00005870
0472      8280 PRINT 8300,IBUB      00005880
0473      8300 FORMAT ('0', 'OPTIMAL VALUE OF OBJECTIVE FUNCTION:', I15//)      00005890
0474      GO TO 8500      00005900
0475      8350 PRINT 8400      00005910
0476      8400 FORMAT ('0', ' PROBLEM DOES NOT HAVE A FEASIBLE SOLUTION',
1          '/1X, ' BECAUSE THE CAPACITY CONSTRA.NTS CANNOT',
2          '/1X, ' BE SATISFIED.',/)      00005920
0477      8500 PRINT 8550      00005930
0478      8550 FORMAT ('0',*****NORMAL END OF JOB*****',/)      00005940
0479      8600 STOP      00005950
0480      END      00005960

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\*Relevant to the TIMET Subroutine



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APPENDIX B

DICTIONARY OF ZIPCAP SYMBOLIC NAMES

<u>SYMBOLIC NAME</u>	<u>DEFINITION</u>
A(I,J)	VARIABLE COST OF ACTIVITY J USING DESIGN I
B(K)	FIXED COST OF FACILITY K
BRO	BRANCHING VARIABLE WITH VALUE 0, I.E., RIGHT-BRANCH VARIABLE X(I,J)
BRL	BRANCHING VARIABLE WITH VALUE 1, I.E., LEFT-BRANCH VARIABLE X(I,J)
BSOLX(J)	CURRENT BEST SOLUTION CONTAINING X(I,J) VARIABLES
BSOLY(K)	CURRENT BEST SOLUTION CONTAINING Y(K) VARIABLES
BSUM	$B(K) * D(I, J, K) / S(K)$ SUMMED OVER APPLICABLE K
BUB	BEST UPPER BOUND
BUBS	BEST UPPER BOUND IF SOLUTION MAY BE SUBOPTIMAL WITHIN EPS. $BUBS = BUB / (1 + EPS)$
C(I,J)	COST ELEMENTS FOR THE RELAXED PROBLEM
CX(I,J)	CONTAINS FIXED AND FREE X(I,J) VARIABLES. FIXED VARIABLES ARE ASSIGNED A VALUE OF 1 OR 2, AND FREE VARIABLES ARE ASSIGNED A VALUE 0. A VALUE OF 1 IMPLIES THAT THAT PARTICULAR X(I,J) HAS A VALUE 1, AND FIX(J) IS FIXED AS 1. A VALUE OF 2 IMPLIES THAT THAT PARTICULAR VARIABLE SHOULD NOT BE CONSIDERED FOR CURRENT COMPUTATIONS CX(I,J) IS UPDATED BY THE CAPACITY RULE, THE BOUNDING RULE, AND THE RULES FOR BRANCHING AND BACKTRACKING
D(I,J,K)	CAPACITY REQUIRED AT FACILITY K FOR ACTIVITY J WHEN ACTIVITY J USES DESIGN I
DBOUND	DIFFERENCE BETWEEN BOUNDS BUBS AND LOWB
DIFBR(J)	DIFFERENCE IN C(I,J) VALUES FOR SELECTING BRANCHING VARIABLE
E(I,K)	HAS VALUE 1 IF DESIGN I USES FACILITY K, OTHERWISE THE VALUE IS 0
ELT1	ELAPSED TIME IN SECONDS TO EXECUTE THE PROGRAM OBTAINED FROM THE TIMET SUBROUTINE

ELTN	ELAPSED TIME IN SECONDS TO COMPARE WITH ET IF SPECIFIED
EPS	EPSILON VALUE IF A SUBOPTIMAL SOLUTION ACCEPTABLE, E.G., EPS OF 0.002 IMPLIES THAT THE SOLUTION MAY BE SUBOPTIMAL, BUT IS GUARANTEED TO HAVE A VALUE WITHIN 0.2% OF THE OPTIMAL VALUE. EPS IS 0 IF AN OPTIMAL SOLUTION IS DESIRED
ET	ELAPSED TIME, IF SPECIFIED, AT WHICH THE NODE AND BOUNDS INFORMATION IS PRINTED
FC	FIXED COST FOR COMPUTING LOWER BOUND
FCUB	FIXED COST FOR COMPUTING UPPER BOUND
FIX(J)	VECTOR OF FIXED COLUMNS. A COLUMN IS FIXED IF IT HAS A FIXED X(I,J) VARIABLE WITH VALUE 1
FIXI(J)	VECTOR SPECIFYING INDEX I CORRESPONDING TO FIX(J) OF VALUE 1
FLB(K)	VECTOR OF FACILITIES FOR COMPUTING LOWER BOUND
FUB(K)	VECTOR OF FACILITIES FOR COMPUTING UPPER BOUND
I	INDEX FOR DESIGNS
IBALD	BALANCE AVAILABLE CAPACITY FOR A GIVEN FACILITY
IBNOD	NODE CORRESPONDING TO THE BEST UPPER BOUND
IBUB	BEST UPPER BOUND
ICAPR	OPTION TO USE CAPACITY RULE
IDIFD	DIFFERENCE BETWEEN A D(I,J,K) VALUE AND MIND(J)
IINPT	OPTION TO LIST INPUT DATA
IMIN	MINIMUM OVER INDEX I FOR A GIVEN COLUMN J
INET	INDICATOR FOR ET
INIS	HAS THE SAME VALUE AS ISTEP
INS	INDEX FOR STACK OF VARIABLES STX
INSET	INDEX FOR DISPLAYING DETAILED STEPS IF OPTION ET IS SPECIFIED
INT	INTERMEDIATE CALL TO TIMET SUBROUTINE

ISTEP	OPTION TO LIST SUMMARY OF INTERMEDIATE STEPS, DETAILED STEPS, OR TO SKIP THE LISTING
ITO	INITIAL CALL TO TIMET SUBROUTINE
IT1	FINAL CALL TO TIMET SUBROUTINE
IUNCAP	OPTION TO INDICATE IF THE PROBLEM BEING SOLVED IS CAPACITATED OR UNCAPACITATED
IX	INDEX I OF VARIABLE $X(I,J)$
J	INDEX FOR ACTIVITIES
JX	INDEX J OF VARIABLE $X(I,J)$
K	INDEX FOR FACILITIES
KT2(J)	A COUNTER FOR $CX(I,J)$ VALUES SET EQUAL TO 2 FOR COLUMN J
LK	A COUNTER
LOWB	LOWER BOUND
LQ1	A COUNTER FOR $CX(I,J)$ VALUES SET EQUAL TO 1
LQ2	A COUNTER FOR $CX(I,J)$ VALUES SET EQUAL TO 2 IN A CYCLE IN WHICH ALL THE FACILITIES ARE EXAMINED
LR2	A COUNTER SET EQUAL TO LQ2 AT THE END OF A CYCLE WHEN APPLYING THE CAPACITY RULE
P	NUMBER OF FACILITIES
S(K)	TOTAL CAPACITY OF FACILITY K
SOLX(J)	SOLUTION SHOWING INDEX 1 OF $X(I,J)$ VARIABLES CORRESPONDING TO $J=1, 2, \dots, N$
STX(INS)	STACK OF FIXED $X(I,J)$ VARIABLES FOR BRANCHING AND BACK-TRACKING. AN $X(I,J)$ RECORDED IN $CX(I,J)$ AS 1 DUE TO THE BRANCHING RULE IS RECORDED IN STX(INS) AS $X * 100 + J$ . AN $X(I,J)$ RECORDED IN $CX(I,J)$ AS 1 DUE TO THE CAPACITY RULE OR THE BOUNDING RULE IS RECORDED IN STX(INS) AS $(X * 100 + J) + 1,000,000$ . AN $X(I,J)$ RECORDED IN $CX(I,J)$ AS 2 IS RECORDED IN STX(INS) AS $-(X * 100 + J) - 1,000,000$
TIMET	A SUBROUTINE FROM THE PL1 LIBRARY TO RECORD THE EXECUTION TIME

UPB

UPPER BOUND

ZIPCAP

NAME OF THE PROGRAM, AN ACRONYM FOR ZERO-ONE INTEGER  
PROGRAM TO SOLVE MULTIACTIVITY MULTIFACILITY CAPACITY-  
CONSTRAINED ASSIGNMENT PROBLEMS

T-423

APPENDIX C  
ANNOTATED OUTPUT FOR TEST PROBLEM 1

OPTIONS SELECTED : IJNPI=1 ICAPR=1 ISTEP=1 IUNCAP=0 EPS= 0.0 ET= 0.0

INPUT DATA

NUMBER OF DESIGNS (M)= 3  
NUMBER OF ACTIVITIES (N)= 4  
NUMBER OF FACILITIES (P)= 5

VARIABLE COST MATRIX A(I,J)

I= 1	9847	15718	4450	8925
I= 2	10192	10542	4422	8118
I= 3	11019	9750	4409	8337

FIXED COST VECTOR B(K)

1750	2000	1750	1350	1000
------	------	------	------	------

CAPACITY LIMIT VECTOR S(K)

400	400	1000	400	400
-----	-----	------	-----	-----

CAPACITY USAGE MATRIX D(I,J,K)

K= 1				
I= 1	0	0	0	0
I= 2	186	60	261	148
I= 3	154	38	175	100

③

②

K= 2					
I= 1	0	0	0	0	0
I= 2	0	0	0	0	0
I= 3	154	38	175	100	100
K= 3					
I= 1	525	144	1011	396	396
I= 2	186	60	261	148	148
I= 3	154	38	175	100	100
K= 4					
I= 1	0	0	0	0	0
I= 2	0	0	0	0	0
I= 3	154	38	175	100	100
K= 5					
I= 1	0	0	0	0	0
I= 2	186	60	261	148	148
I= 3	154	38	175	100	100
<b>DESIGN-FACILITY MATRIX E(I,K)</b>					
I= 1	0	0	1	0	0
I= 2	1	0	1	0	1
I= 3	1	1	1	1	1

## DESIGN=FACILITY MATRIX E(I,K)

0	0	1	0	0	1
1	0	1	0	1	2
1	1	1	1	1	3

## SUMMARY OF STEPS

0 37229-3594

37229.3594

**MINSC, FC, LOW**

for the



## NODE NUMBER 2

MINSC, FC, LOWB 35528.3633 1750 37278.3633  
 BR1 203

## NODE NUMBER 3

MINSC, FC, LOWB 32356.0000 7850 40206.0000  
 NSUMA, FCUB, UPB, BUB, BUWS 32356 7850 40206.0000 999999.0000 999999.0000  
 BRO 203

## NODE NUMBER 4

MINSC, FC, LOWB 32324.0000 7850 40174.0000  
 NSUMA, FCUB, UPB, BUB, BUWS 32324 7850 40174.0000 40206.0000 40206.0000  
 BRO 101

## NODE NUMBER 5

MINSC, FC, LOWB 33476.0000 7850 41326.0000

ELAPSED TIME IN SECONDS= 0.06226163

TOTAL NUMBER OF NODES EXPLORED = 5

NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.  
 2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

## OPTIMAL SOLUTION

X(I,J) WITH VALUE 1: 1 3 3 2  
 Y(I) VALUES: 1 1 1 1 1  
 OPTIMAL VALUE OF OBJECTIVE FUNCTION: 40174

\*\*\*\*\*NORMAL END OF JOB\*\*\*\*\*

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APPENDIX D  
ANNOTATED OUTPUT FOR TEST PROBLEM 2

OPTIONS SELECTED : IINPT=1 ICAPR=1 ISTEP=0 IUNCAP=0 EPS= 0.00200 ET= 0.0

INPUT DATA

NUMBER OF DESIGNS (M)= 5  
 NUMBER OF ACTIVITIES (N)= 4  
 NUMBER OF FACILITIES (P)= 8

VARIABLE COST MATRIX A(I,J)

I= 1	194951	218871	155094	104056
I= 2	235277	272087	143264	138641
I= 3	196138	220718	160399	107481
I= 4	198751	224042	167046	112445
I= 5	190873	229169	128361	112498

FIXED COST VECTOR B(K)

14000	14000	14000	14000	14000	14000	28000	19000	31000
-------	-------	-------	-------	-------	-------	-------	-------	-------

CAPACITY LIMIT VECTOR S(K)

350	350	350	350	200	700	500	800
-----	-----	-----	-----	-----	-----	-----	-----

CAPACITY USAGE MATRIX D(I,J,K)

K= 1

©

1= 1	80	90	100	180
1= 2	0	0	0	0
1= 3	80	90	100	180
1= 4	40	30	50	120
1= 5	0	0	0	0
1= 2				
1= 1	80	90	100	180
1= 2	0	0	0	0
1= 3	80	90	100	180
1= 4	40	30	50	120
1= 5	0	0	0	0
1= 3				
1= 1	80	90	100	180
1= 2	0	0	0	0
1= 3	80	90	100	180
1= 4	40	30	50	120
1= 5	0	0	0	0
1= 4				
1= 1	80	90	100	180
1= 2	0	0	0	0
1= 3	80	90	100	180
1= 4	40	30	50	120
1= 5	80	90	100	180
1= 5				
1= 1	80	90	100	180
1= 2	0	0	0	0
1= 3				

(c)

I= 4	40	30	50	120
I= 5	80	90	100	180
K= 6				
I= 1	240	180	300	450
I= 2	0	0	0	0
I= 3	120	90	150	360
I= 4	0	0	0	0
I= 5	40	30	50	120
K= 7				
I= 1	160	180	200	360
I= 2	0	0	0	0
I= 3	80	60	100	240
I= 4	0	0	0	0
I= 5	160	180	200	360
K= 8				
I= 1	200	150	250	600
I= 2	40	30	50	180
I= 3	0	0	0	0
I= 4	0	0	0	0
I= 5	120	90	150	300

DESIGN-FACILITY MATRIX E(I,K)

I= 1	1	1	1	1	1
I= 2	0	0	0	0	1
I= 3	1	1	1	1	0
I= 4	1	1	1	0	0

I= 5                    0                    0                    0                    1                    1                    1                    1                    1                    1

ELAPSED TIME IN SECONDS= 0.09392625

TOTAL NUMBER OF NODES EXPLORED = 11

NOTE: 1. FOLLOWING X(I,J) VARIABLES SHOW DESIGN I TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.  
2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

OPTIMAL SOLUTION

X(I,J) WITH VALUE 1:                    4                    4                    2                    4  
Y(I) VALUES:                    1                    1                    1                    1                    1                    0                    0                    1  
OPTIMAL VALUE OF OBJECTIVE FUNCTION: 779502

\*\*\*\*\*NORMAL END OF JOB\*\*\*\*\*

T-423

APPENDIX E  
ANNOTATED OUTPUT FOR TEST PROBLEM 3

OPTIONS SELECTED : IIMPT=1 ICAPR=0 ISTEP=0 IUNCAP=1 EPS= 0.0 ET= 30.000

INPUT DATA

NUMBER OF DESIGNS (M)= 10  
NUMBER OF ACTIVITIES (N)= 8  
NUMBER OF FACILITIES (P)= 8

VARIABLE COST MATRIX A(I,J)

I= 1	21	13	20	0	15	10	35	20
I= 2	26	13	27	14	17	13	40	25
I= 3	25	12	26	13	16	12	20	23
I= 4	23	9	22	20	19	20	50	19
I= 5	23	20	24	13	20	15	40	17
I= 6	23	20	24	13	20	15	40	17
I= 7	26	13	24	20	15	15	35	20
I= 8	30	20	40	23	30	20	55	25
I= 9	26	13	26	20	20	15	40	15
I= 10	34	46	44	50	19	20	40	20

FIXED COST VECTOR B(I)

10	8	4	10	14	30	20	50
----	---	---	----	----	----	----	----



①

[illegible]





②

L= 2	1	1	1	1	1	1	1
L= 3	1	1	1	1	1	1	1
L= 4	1	1	1	1	1	1	1
L= 5	1	1	1	1	1	1	1
L= 6	1	1	1	1	1	1	1
L= 7	0	0	0	0	0	0	0
L= 8	0	0	0	0	0	0	0
L= 9	1	1	1	1	1	1	1
L=10	0	0	0	0	0	0	0
L= 7							
L= 1	1	1	1	1	1	1	1
L= 2	0	0	0	0	0	0	0
L= 3	0	0	0	0	0	0	0
L= 4	1	1	1	1	1	1	1
L= 5	1	1	1	1	1	1	1
L= 6	0	0	0	0	0	0	0
L= 7	0	0	0	0	0	0	0
L= 8	0	0	0	0	0	0	0
L= 9	1	1	1	1	1	1	1
L=10	0	0	0	0	0	0	0
L= 8							



ELAPSED TIME IN SECONDS= 0.57446837

TOTAL NUMBER OF MODES EXPLORED = 35

NOTE: 1. FOLLOWING  $X(I,J)$  VARIABLES SHOW DESIGN 1 TO WHICH ACTIVITY J IS ASSIGNED FOR J=1 TO N.

2. IF EPSILON EPS WAS ASSIGNED A POSITIVE (NON-ZERO) VALUE, THE SOLUTION MAY BE SUBOPTIMAL.

OPTIMAL SOLUTION

$X(I,J)$  WITH VALUE 1: 6 3 6 3 3 3 3 6

$V(I,K)$  VALUES: 1 1 1 0 1 1 0 1

OPTIMAL VALUE OF OBJECTIVE FUNCTION: 253

\*\*\*\*\*NORMAL END OF JOB\*\*\*\*\*

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